

1 Introduction

We introduce and model a new class of two-sided matching markets without explicit transfers, in which there is an additional fundamental constraint.¹ The eventual market outcome is linked to an initial status-quo matching, which may give participants certain rights that constrain how future activity can play out. Since market outcome is typically different from the status quo, such activities loosely resemble an exchange in which one side of the agents are changing or acquiring new partners in addition to the two-sided matching market structure. In such markets, a fundamental balancedness condition needs to be sustained with respect to the status-quo matching. The motivation for such a balancedness constraint can be different depending on the features of the market. Two contrasting examples are labor and higher education markets, where workers and colleges provide services to be compensated, respectively. In worker exchange, a worker needs to be replaced with a new one at her home firm so that this firm can function properly, and thus, the market needs to clear in a balanced manner. In student exchange, the college that is matched with an exchange student should be able to send out a student as well so that its education costs do not increase, and thus, the market needs to clear in a balanced manner. There are several prominent examples of such exchanges, such as national and international teacher-exchange programs, clinical-exchange programs for medical doctors, worker-exchange programs within or across firms, and student-exchange programs among colleges. This balancedness constraint induces preferences for firms/colleges not only over whom they get matched with (i.e., import), but also over whom they send out (i.e., export). The most basic kind of such preferences requires the firm/college to have a preference for balanced matchings, i.e., for import and export numbers to be equal. We analyze our model over two explicit market applications: (permanent) tuition exchange and temporary worker exchange (see Section 2 for details).

In tuition exchange

their preferences over matchings are determined through their rankings over the incoming class and how balanced the eventual matching is.² We start by showing, through a simple example, that individual rationality and nonwastefulness, standard concepts in two-sided matching markets, and balancedness are in general conflicting requirements (Proposition 1). For this reason, we restrict our attention to the set of balanced-efficient mechanisms. Unfortunately, there exists no balanced-efficient and individually rational mechanism that is immune to preference manipulation for colleges (Theorem 2).

Although 2S-TTC is balanced-efficient, it may not match the maximum possible number of students while maintaining balance. We show that if the maximal-balanced solution is different from the 2S-TTC outcome for some preference profile, it can be manipulated

Compared to deferred-acceptance-based current practice, they show that a TTC-based approach doubles the number of teachers moving from their initial assignment. Additionally, when the distribution of the ranks of teachers over the schools are considered, the outcome of the TTC-based approach stochastically dominates that of the current practice. Thus, there exist real-life settings, in which our proposals can lead to significant welfare improvements.

We extend this model for temporary worker exchanges, such as teacher-exchange programs. We tweak our model slightly and assume that the quotas of the firms are fixed at the number of their current employees, and, hence, firms would like to replace each agent who leaves. We also assume that firm preferences are coarser than colleges in tuition exchange due to the temporary nature of the exchanges. We assume they have weakly size-monotonic preferences over workers: larger groups of acceptable workers are weakly better than weakly smaller groups of acceptable workers when the balance of the matching with larger groups of acceptable workers is zero and the balance of the matching with smaller group of worker is nonpositive.⁹

Many colleges give qualified dependents of faculty tuition waivers. Through a tuition-exchange program, they can use these waivers at other colleges and attend these colleges for free. The dependent must be admitted to the other college

rent form.¹⁸ The Jesuit universities exchange program FACHEX is another one that is adversely affected. The program still does not have an explicitly embedded balancedness requirement. It includes all Jesuit universities but Georg

students by c . Let $\succ_c = (\succ_c)_{c \in C}$ be the list of college **internal priority orders**, where \succ_c is a linear order over S_c .

In a revelation game, students and colleges report their preferences; additionally, colleges report their admission and eligibility quotas.²⁶ A mechanism is **immune to preference manipulation for students** or **colleges** if for all $[q, e, \succ]$, there exists no $i \in S$ (or $i \in C$) and \succ_i such that $[q, e, (\succ_i, \succ_{-i})](i) \succ_i [q, e, \succ](i)$. A mechanism is **immune to preference manipulation** if it is immune to preference manipulation for both students and colleges. A mechanism is **immune to quota manipulation** if for all $[q, e, \succ]$, there exists no $c \in C$ and (q_c, e_c) with $q_c \neq q_c$ such that $[(q_c, q_{-c}), (e_c, e_{-c}), \succ](c) \succ_c [q, e, \succ](c)$. A mechanism is **strategy proof for colleges** if for all $[q, e, \succ]$, there exists no $c \in C$ and (q_c, e_c, \succ_c) with $q_c \neq q_c$ such that $[(q_c, q_{-c}), (e_c, e_{-c}), (\succ_c, \succ_{-c})](c) \succ_c [q, e, \succ](c)$. A mechanism is **strategy proof for students** if it is immune to preference manipulation for students. A mechanism is **strategy proof** if it is strategy-proof for both colleges and students.²⁷ A mechanism is **group strategy proof for students** if for all $[q, e, \succ]$, there exists no $S \subseteq S$ and

Throughout our analysis, we impose a weak restriction on college preferences. Assumption 1 below states that a college prefers a better scholarship class with zero net balance to an inferior scholarship class with a nonpositive net balance.

Assumption For any μ, M^u and $c \in C$, if $b_c^\mu = 0$, $b_c \leq 0$, and $\mu(c)P^c$

2017). In contrast, in our market, college slots are not objects. Therefore our definition of a mechanism, and the properties of matchings and mechanisms (except strategy-proofness for students) do not have any analogous translation in such problems. However, because

Theorem Under Assumption 1 and when true eligibility quotas satisfy $e_c = |S_c|$ for all $c \in C$, 2S-TTC is immune to quota manipulation.

We prove the theorem with a lemma showing that as the quotas of a college increase, the import and export sets and the admitted class of students of this college also (weakly) expand under 2S-TTC.³²

Theorems 3 and 4 point out that only colleges can benefit from manipulation, and they can manipulate by misreporting their preferences. Moreover, the only way to manipulate preferences is to report an acceptable student as unacceptable. Suppose we take all the admitted students in the regular admission procedure as acceptable for a tuition-exchange scholarship. Then, to manipulate 2S-TTC, a college needs to reject a student who satisfies the college admission requirements. Usually college admission decisions are made before the applicants are considered for scholarships.³³

Proposition 2 below implies that colleges do not benefit from misreporting their ranking over incoming classes.

Proposition Under Assumption 1, colleges are indifferent among strategies that report preferences in which the same set of students is acceptable with the same quota report under the 2S-TTC mechanism.

We have shown that 2S-TTC has appealing properties. In the following theorem, we show that it is the unique mechanism satisfying a subset of these properties.

Theorem 5 Under Assumption 1, 2S-TTC is the unique student-strategy-proof, acceptable, and balanced-efficient mechanism that also respects internal priorities.

In the proof of our characterization theorem, we use a different technique from what is usually employed in elegant single quota characterization proofs such as Svensson (1999) and Sönmez (1995) for the result of Ma (1994). Our proof relies on building a contradiction with the claim that another mechanism with the four properties in the theorem's

³² For instance, contrast with the results in the literature on stable matchings, student and college optimal stable matchings are prone to a loss on quota manipulation by the college. Svensson (1999) provides a proof of this result. For an alternative proof, see Konrad and Lommerud (2008). C

hypothesis can exist. Suppose such a mechanism exists and finds a different matching than 2S-TTC for some market. The 2S-TTC algorithm runs in rounds in which trading cycles are constructed and removed. Suppose $S(k)$ is the set of students removed in Round k , while running 2S-TTC in such a way that in each round only one arbitrarily chosen cycle is removed and all other cycles are kept intact. We find a Round k and construct an auxiliary market with the following three properties: (1) Eligibility quotas of home colleges of students in $S(k)$ are set such that these are the last certified students in their respective home institutions; (2) all preferences are kept intact except those of students in $S(k)$, whose preferences are truncated after their 2S-TTC assignments; and (3) all students in $S(k)$ are assigned c under the alternative mechanism, while all students removed in the 2S-TTC algorithm before Round k have the same assignment under 2S-TTC and the alternative mechanism. This contradicts the balanced-efficiency of the alternative mechanism: we could give the students in $S(k)$ their 2S-TTC assignments while keeping all other assignments intact and obtain a Pareto-dominating balanced matching. Round k and the auxiliary market are constructed in three iterative steps.

Among all the axioms, only the respect for internal priorities is based on exogenous rules. One might suspect that more students will benefit from the tuition-exchange program if we allow the violation of respect for internal priorities. However, such mechanisms turn out to be manipulable by students.

Theorem 6 Any balanced and individually rational mechanism that does not assign fewer students than 2S-TTC and selects a matching in which more students are assigned whenever such a balanced and individually rational outcome exists, is not strategy-proof for students, even under Assumption 1.

4.1 Market Implementation: Tuition Remission and Exchange

Incorporating tuition-remission

in a semi-decentralized fashion: first, colleges announce their tuition-exchange scholarship quotas and which of their students are eligible to be sponsored for both exchange and remission; then, eligible students apply for scholarship to the colleges they find acceptable; then colleges send out scholarship admission letters. At this stage, as students have also learned their opportunities in the parallel-running regular college admissions market, they can form better opinions about the relative ranking of the null college, i.e., their options outside the tuition-exchange market. Students submit rankings over the colleges that admitted them with a tuition-exchange scholarship and the relative ranking of their outside option. Finally, 2S-TTC is run centrally to determine the final allocation.

4.2 Allowing Tolerable Imbalances

Some programs care about approximate balance over a moving time window. Here, we relax the zero-balance constraint and allow each $c \in C$ to maintain a balance within an interval $[l_c, u_c]$ where $l_c \leq 0 \leq u_c$.³⁴ When either l_c or u_c equals zero for all $c \in C$, the market turns into the case studied in Section 4. Let $(l_c, u_c)_{c \in C}$ be the tolerance profile.

When the colleges hold a non-zero balance, then there may exist some colleges exporting (importing) more than they import (export). Then, we cannot represent all allocations by cycles. Therefore we need to consider chains in addition to the cycles. A

which consider her acceptable, and c

Theorems 7 and 8 hold without any assumptions on preferences. Under a mild assumption on college preferences, we can show that 2S-TTTC is individually rational and it induces a dominant-strategy equilibrium for colleges' quota reporting game to certify all their students and report their true admission quota.

Although 2S-TTTC is defined in a static problem, we can easily extend it to the dynamic environment where the aggregate balance over years matters. In particular, for each period t and $c \in C$ we can set counter b_c equal to c 's aggregate balance in period $t - 1$ where the aggregate balance in period $t - 1$, is equal to the sum of balances between period 1 and $t - 1$. Moreover, the exogenous priority rule used in period t can be determined based on the aggregate balance colleges carry at the end of period $t - 1$ such that the highest priority can be given to the college with the highest aggregate balance and so on.

5 Temporary Worker Exchanges

Many organizations have temporary worker-exchange programs that can be modeled through our balanced two-sided matching framework. The first difference between such programs and tuition exchange is that these exchanges are usually temporary. Each firm usually requires a set of specific skills, e.g., a mathematics teacher to replace their own mathematics teacher. Compatibility and ability to perform the task are the main preference criterion rather than a strict preference ranking. E.g., finding a good teacher with a specific degree is the first-order requirement, rather than finer details about the rankings of all good teachers.

The second difference is that each position and each worker should be matched, unlike the tuition-exchange application. The workers are currently working for their home firms. Thus, the firms consider these workers necessarily acceptable. By contrast, in tuition exchange, colleges are not required to admit all the dependents of their employees. In temporary worker exchanges, a worker who does not want to go to a different firm necessarily stays employed in her home firm. We need to use a variant of the tuition-exchange model to facilitate balanced-efficient trade in such circumstances.

We can use the model introduced in Section 3 with slight changes. Since each firm accommodates its current workers, $q_c = |S_c|$ for each $c \in C$ if $g_T(c) = 919955$ # e (c) 9199

Our paper, besides introducing a new applied problem and proposing a solution to it, has six main theoretical and conceptual contributions: We introduce a new two-sided matching model that builds on the two most commonly used matching models in the literature: discrete object allocation, including school

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Supplementary Material for
"Two-Sided Matching via Balanced Exchange"
by
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Appendix A On Current Practice of Tuition Exchange

In this appendix we analyze the current practice of tuition exchange. As the centralized process is loosely controlled, once each college sets its e.g. b.ty admission quota and e.g. b.e students are determined, the agent functions for

assumption states that a better admitted class is preferable as long as the net balance does not decrease, admission of unacceptable students deteriorates the rankings of unconstrained applicants regardless of their net balances, and a college desires its own students unacceptable tuition exchange. Assumption 4 introduces negative net balance averse preferences. In a results section we will use Assu

We prove this proposition by constructing an associated Gale \mathbb{D} -decomposable admissions
and then with the set of Gale \mathbb{D} -decomposable allocations.

countries and do not conduct an equilibrium analysis in a quota determination game. But they do point out that in a frictionless market the countries that would be likely to have a negative balance would be conservative and would decrease their export quotas for exports which would further deteriorate the balances of other countries.

Typically no country would draw in practice as there is often a national quota of participation in place. The conjecture that this could be instituted because of the reasons outlined above. Given that continued membership is an attractive benefit often times a country would announce that they would import and export at this national quota requirement and would continue to be a member of the program without further drawing from the system.

We conclude that under a new design for tuition exchange there should be no room

Acceptability: Students will be assigned to null college c_0 whenever they point to t and hence they will never need to point to an unacceptable college. Hence a student cannot be assigned to an unacceptable college. Moreover, a student cannot point to a college that considers her unacceptable therefore the students ranked below n_c cannot be assigned to c unless c is acceptable.

Individual Rationality: Once each student is assigned to an option weakly better than c_0 she does not find a strictly better option. Once a student $n_c(c)$ are ranked above n_c for each c .

students who are not in the cycles removed in round K of \mathcal{D} . C were K stable in the last round of \mathcal{D} . C^{10} we prove that μ is balanced efficient in two parts

Part I: we first prove that μ cannot be Pareto dominated by another acceptable balanced matching. If $s \in \mu^{-1}(c)$ then $(s, c) \in C$ is the highest ranked option in μ that s considers acceptable at s , so no student $s' \in \mu^{-1}(c)$ can be assigned to a better college considering μ acceptable. If there exists a matching μ' such that $\mu' \succ s$, then (s, c) considers s unacceptable at s , so μ cannot be Pareto dominated by another acceptable matching in $w.c.$ at least one student $s' \in \mu^{-1}(c)$ is better off in

If a student $s' \in \mu^{-1}(c)$ is not assigned to a more preferred $c \in C$ that s' considers acceptable then c should be removed in round 1. Let μ be an acceptable and balanced matching such that $\mu(s) = c$. Suppose there exists another student s' such that $\mu(s') = c$ and $(s', c) \in C$. Note that s' is an eligible student. Because s' is assigned in round 1, (s', c) is her favorite college among the ones considering μ acceptable at s' in any acceptable and balanced matching in $w.c.$ so s' is assigned to (s') , s' will be made worse off. Suppose $(s', c) \in C$ for any $s' \in \mu^{-1}(c)$. ¹¹ $\mu \in C$

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Lemma 1

Proof of Theorem 4. We prove a stronger version of Theorem 4 under $\mathcal{D} \subseteq \mathcal{C}$. Suppose that preference profiles are fixed for colleges such that no college reports an unacceptable student as acceptable in its preference report. In the induced quota reporting game under Assumption 1, it is a dominant strategy equilibrium for a college $c \in \mathcal{C}$ to certify

en any acceptable ec an s w ass gn er to e_0 If $|(1)| = 1$ or $|(1)| = 1$

one cycle

Step 1: Construct a preference profile $\tilde{\lambda}$ with associated ranking \tilde{r} as follows. Let student s have ranking $r_c(s)$ as acceptable preferences and $\tilde{\lambda}_j = \lambda_j$ for all $j \in [(C-1) \setminus s]$. By the execution of the Gale-Shapley algorithm, μ is a stable matching for $[q, \lambda]$. Since μ is a strategy proof for students and acceptable preferences $[q, \tilde{\lambda}](s) = c_\emptyset$

then we can get either the assignments of students in $K_{k'=1}^c(\mu)$ are the same in $[q, \tilde{\lambda}]$ and μ . If not then for some $k \in K_{k'=1}^c(\mu)$, there exists a student $\tilde{s} \in K_{k'=1}^c(\tilde{\mu})$ preferring \tilde{s} to $\mu(k)$ and each student in $K_{k'=1}^c(\mu)$ gets the same college in $[q, \tilde{\lambda}]$. Then we repeat Step 1 by taking $\tilde{\lambda} := \tilde{\lambda}$, $s := \tilde{s}$ and $\mu := \tilde{\mu}$.

This repetition will end by the finiteness of rounds. Then all students in $K_{k'=1}^c(\mu)$ get the same college in $[q, \tilde{\lambda}]$. The Gale-Shapley algorithm outputs $[q, \tilde{\lambda}]$ and $[q, \tilde{\lambda}]$ then we proceed to Step 2.

Step 2: In Step 1 we have shown that s prefers $\mu(s)$ to $[q, \tilde{\lambda}](s) = c_\emptyset$. Suppose c is the college of s . Set a new eligibility quota \tilde{c}_c equal to the ranking of student s in c in its internal priority order that is $\tilde{c}_c = r_c(s)$ and let

colleges which consider her acceptable and c_0 . If a student is assigned in this round then she should get the same college. Now consider students assigned in round t when $t > 1$. At the colleges that a student prefers to her assignment and considers acceptable should have been removed or become non-reporting in an earlier round. She cannot get this student better off by assigning her to a college that considers her

Once we can run \mathcal{D} on a given assignment of workers to territories, the proof of the existence of a worker strategy proofness of \mathcal{D} on the proof of the existence of a stable assignment. First note that any Pareto efficient student strategy proof and acceptable mechanism assigns workers to territories or better results at the acceptable. In the uniqueness part of the proof, the mechanism's proof adopted for \mathcal{D} being the only efficient satisfying Pareto efficiency student strategy proofness, acceptability, and respect for internal priorities in the temporary worker exchange mode, where updating workers' preferences in Step 1, we do it as follows: rank (s) and territories as only acceptable results in the correct order of their true preferences. And then at the end of Step 1, s will be assigned to territories under \succ . Once \succ respects internal priorities and is acceptable student strategy proof and balanced efficient, s will be an efficient result in Step 2. Then we reach Step 3, we will have a set of workers who are assigned to territories by \succ . However, a trading cycle between the world will prove that we fare with out violating balancedness or feasibility.

Immunity to Preference Manipulation by Colleges: Because at any allocation, balancedness is satisfied and results are efficient on quotas. Hence, under Assumption 1, results are indifferent between any acceptable allocation. Once the mechanism selects an acceptable allocation when results report truthfully, results cannot be better off by manipulating their preferences over the allocations and reporting quotas different from their true quotas.

Stability: Consider an arbitrary market $[q, \succ]$. Denote the outcome of \mathcal{D} on C by (c) at $q_c = |c|$ for a given $C \subseteq C$. Workers consider their current results acceptable, results consider their current workers acceptable, and workers who are not certified are an efficient result. Hence, $|c| = q_c$ for a given $C \subseteq C$. Once in a given results quotas are efficient, nonwasteful. Note that any mutual deviation of workers' reports needs to end up with a balanced allocation. Once a given employees in (c) are acceptable, replacing one of the employees in (c) with another one in $C \setminus (c)$ cannot be better off. Hence, (c) cannot be blocked by a worker's report. ■

Appendix C Tuition-Exchange Programs

The first explanation why tuition exchange programs exist in the first place is because some colleges choose to subsidize faculty directly instead of participating in tuition exchange programs. Although this may create a benefit for the students, any direct compensation

over $\frac{1}{2}$ is taxable income whereas a tuition exchange scholarship is not²³ tuition exchange is not considered to be an income transfer²⁴ Moreover colleges may not want to switch to successful compensation programs from a cost saving perspective regardless of the tax benefit to the faculty member we present a simple back of the envelope calculation to demonstrate these cost savings there are more than 1000 year colleges in the US and at most a fraction of the average members participate in at least one tuition exchange program Suppose n students are given tuition exchange releases on scholarships a year Instead of a college finances the tuition of a faculty member's child through direct cash compensation then a tuition exchange colleges would have to pay $\$n\bar{c}$ where \bar{c} is the average future tuition cost of colleges However assuming that average quantities and sizes of colleges with and without tuition scholarships are the same only a fraction of these students would attend a tuition exchange college in return so the colleges would only get back $\frac{n\bar{c}}{2}$ the remaining $\frac{n}{2}$ slots would be filled with regular students regular students on average pay about a fraction of the tuition than to other financial programs For example a tuition discounting Study of the National Association of College and University Business Officers report that incoming freshmen pay on average 56% of future tuition at a private university thus they would only pay $\frac{n\bar{c}}{4}$ to tuition exchange colleges As tuition exchange scholarships constitute a very small portion of college admissions this calculation assumes that average tuition payment would not change by establishment of direct cash compensation instead of tuition exchange thus as a result the colleges would lose in total about $\frac{n\bar{c}}{4}$ which corresponds to one fourth of average future tuition per student thus the total per student savings for the faculty member and the colleges are more than a fraction of tuition payment assuming one third of the direct compensation is paid in income tax at the margin by the parent

The Tuition Exchange Inc (TTEI):3(a92(3(6(m)3(e)4(n)3g)120011.955297.1991440.399

year in each faculty member submits the EI application to the registration office of their college. If the number of applicants is greater than the number of students that the college is willing to certify then the college decides who to certify based on years of service or some other criterion in the internal priority order.

Each student who is certified eligible submits a list of colleges to the liaison office of their institution. Each liaison office sends a copy of the EI Certificate of Eligibility to the EI liaison officer at the participating colleges and universities listed by the eligible dependents. Certification only means that the student is eligible for an EI award. It is not a guarantee of an award. The eligible student must apply for admission to the college(s) in which she is interested following each institution's application procedures and deadlines. After admission decisions have been made the admissions offices or EI liaisons at the listed institutions inform her whether she will be offered an EI award.

EI scholarships are competitive and so the eligible applicants may not receive the award. The sponsoring institution cannot guarantee that an export candidate regardless of qualifications will receive an EI scholarship. Institutions choose their scholarship recipients based on the applicants' academic profiles.

To collect anecdotal evidence on how our faculty members value the tuition exchange benefit we also conducted an I/B approved evaluation survey of the tuition exchange colleges and EI members and possibly members of other tuition exchange programs using Qualtrics evaluation survey software. Our respondent pool consisted of 10 faculty members with a 75 to 85% response rate. In this pool there are 2 and 1 assistant and 1 full-time and 1 full-time.

employees e.g. be based on its own rules. Each member college is required to accept at least three exchange students per year. There is no limitation on the number of exported students. Each certified student also applies for admission directly to the member colleges of her choice. Certified students must be admitted by the host college in order to be considered for the tuition exchange scholarship. Each year more than 100 students benefit from this program.

Catholic College Cooperative Tuition Exchange (CCCTE) CCCTE is composed of member colleges. Each member college certifies its employees as e.g. be based on its own rules. Students must be admitted by the host college before applying for

exchange of positions with teachers from countries including the Greece, Finland and Netherlands, India, Mexico and the UK. Matching procedures arranged by the Fulbright program staff, and each candidate and each school must be approved before the placements are finalized.

The Commonwealth Teacher Exchange Programme (CTEP) was founded by the League for the Exchange of Commonwealth Teachers more than 10 years ago. Participant teachers exchange their observations and ideas with each other usually for a year, and they stay employed by their own schools. Countries participating in this program are Australia, Canada, and the UK. More than 100 teachers have benefited from the CTEP. Principals have the right to veto any proposed exchange they think would not be appropriate for their schools.

The Educator Exchange Program is organized by the Canadian Education Exchange Foundation. The program includes reciprocal, interprovincial, and international exchanges. The international destinations are Australia, Denmark, France, Germany, Switzerland, and the UK and Colorado, the US.

The Manitoba Teacher Exchange enables teachers in Manitoba to exchange their positions with teachers in Australia, the UK, Germany and other Canadian provinces. Once a potential placement is found, the incoming teacher's information is sent to the Manitoba applicant, the principal of the school, and the employing authority. Acceptance of all these teachers is required for the completion.

possible exchange counterpart then they can exchange their positions before entering the

The Erasmus Student Exchange Program is a leading exchange program between the universities in Europe. Close to 1 million students have participated since it started in 1987. The number of students benefiting from the program is increasing each year. More than 1 million students attended a college in another European country as

Appendix E Proofs of Appendix A

Proof of Proposition 3. We prove existence by showing that for any tuition exchange market there exists an associated college admissions market and the set of stable matchings are the same under both markets. Under Assumption 1, we fix a tuition exchange market (q, \succsim) . Let E be the set of eligible students. We first introduce an associated college admissions market, a Gale-Debreu-Lyapunov two-sided many-to-one matching market (C, q, s, \succ) where the set of students is E , the set of colleges is C , the quota vector of colleges for admissions is q , the preference profile of students over colleges is \succ which are all the same as reported from the tuition exchange market and the preference profile of colleges over the set of students is \succ .

Finally we show that factoring is not stable for

3.b.i. If $\tilde{c} = c$

and let $(q'_c, \hat{q}_{-c}) \in (\hat{c}, \hat{c}-1)$ where $\hat{c} = \hat{c} - 1$ ■

Appendix F Structure of Stable Matchings

rankings of agents associated with their preferences over outcomes are given as

					a	b	c	d	e									
\succsim_a	\succsim_b	\succsim_c	\succsim_d	\succsim_e	3	5	2	2	2	1	2	3	4	5	6	7	8	9
1	3	6	7	9	4	1	3	3	3	b	b	a	c	b	a	c		c
2	4	5	8		5	6	4	4	8	c	c	c	a	a	b	a	c	
					9	2	9	9	7	c_\emptyset	c_\emptyset	c_\emptyset	c_\emptyset	c_\emptyset	c_\emptyset	c_\emptyset	c_\emptyset	c_\emptyset
					7	7	7	5	5									

Let o_e and o_a be the vectors representing the endowment and assets on counters of colleges respectively. Then we set $o_e = (2 \ 2 \ 2 \ 2 \ 1)$ and $o_a = (2 \ 2 \ 2 \ 1$

Courses a and c are removed

Round 4: the only cycle formed in round 4 is $(c_0, 7)$ therefore 7 is assigned to c_0
Given that we have a trivial cycle including c_0 we only update o_e

are acceptable for colleges. The ranking associated with preference profiles \succsim_S is given as

1	2	3
b	a	b
c	c	a
a	b	c
c_\emptyset	c_\emptyset	c_\emptyset

Let each agent i select the same matching as in μ^C for each agent except the agent $[q = (1, 1, 1) = (1, 1, 1) \succ]$ and for this agent i assigns **1** to c , **2** to a and **3** to b . This mechanism is balanced efficient, acceptable and respects individual preferences. However, it is not student strategy proof because when **1** reports c unacceptable with assignment **1** to b ,

- A balanced-efficient, student-strategy-proof, but not acceptable mechanism that re-

represents the correlation of tastes of students on c' $X(s, c')$ (0, 1) is also an individual standard

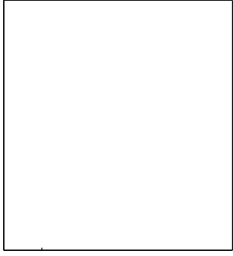
where α is used as the horizontal axis variable and the vertical axis variables in top and bottom graphs demonstrate the difference of the percentage of unassigned students between the DA decision under the two alternative strategies of the colleges. In each row, the left and right graphs are for straightforward behavior of DA, i.e., strategy 1 and the end and start graphs are for the equilibrium behavior of DA, i.e., strategy 2, explained above and in

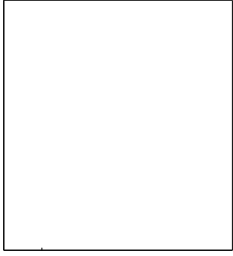
Figure 1. In bottom graphs, the vertical axes demonstrate the difference between the percentage of the students preferring the versions of α and the percentage of the students preferring the DA decision under two alternative strategies of the colleges.³⁴

Under a scenario when we compare the percentage of students preferring the versions of α and the DA decision under two alternative strategies of the colleges, we observe that α and α outperform both alternative strategic behaviors under DA. For example, when $\alpha = 0.5$ and $\alpha = 0.5$ for year y to entrance level, 19.23% more of α students, i.e., the percentage of students who prefer α to DA, minus the percentage who prefer DA to α prefer α outcome to DA straightforward behavior outcome when the test difference increases to

DA behavior scenarios. On the other hand, as the students' preference correlation parameter increases, the dominance measures display mostly a unimodal pattern peaking for moderate ρ for any fixed

and each college has a available seats. Different from the previous cases, the number of students applying to be certified may vary and they are selected from interval $[6, 10]$ according to a uniform distribution. Preference profiles of the students and the colleges are





References

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